# Network Graphs 

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## Overview

(1) Graph
(2) KCL
(3) KVL
(4) Node-based Description
(5) Mesh-based Description
(6) Cut Set-based Description
(7) Loop-based Description
(8) Comparison of Different Descriptions
(9) Duality
(10) Tellegen's Theorem

## Graphs

## Graphs



Figure: Each circuit can be represented by a network graph if each element is replaced with an edge having two ending nodes. The nature of elements is discarded in the network graph. A circuit may have a unconnected graph.

## Graphs

## Definition (graph)

A graph is mathematically described by $G(\boldsymbol{N}, \boldsymbol{E})$, where $\boldsymbol{N}$ is the set of nodes and the set of edges $\boldsymbol{E}=\left\{\left(e_{i}, e_{j}\right) \mid e_{i}, e_{j} \in \boldsymbol{N}\right\}$.


Figure: Graphs with isolated node and self-loop along with a complete graph.

## Graphs



Figure: A graph and some of its subgraphs.

## Graphs



Figure: Associated reference directions for an element and for a branch.


Figure: A network and its corresponding directed graph.

## KCL

## Cut Sets



Figure: Connected and unconnected graphs. A unconnected graph have two or more separated parts.


Figure: Branch removal operation.

## Cut Set

## Definition (Cut Set)

A cut set is the set of branches such that

- The removal of all the branches of the set adds a new separated part to the graph.
- The removal of all but any one of the branches of the set adds no new separated part to the graph.


Figure: Example of cut sets.

## Node and Gaussian surface

## Statement (Node)

A node is a special cut set that only surrounds a node.

## Statement (Gaussian surface)

A Gaussian surface is a generalized cut set that decomposes the graph into two or more separated parts.


Figure: Examples of node and Gaussian surface.

## KCL

## Definition (KCL)

For any lumped network and at any time, the algebraic sum of all the branch currents entering (exiting) a cut set (node, Gaussian surface) branches is zero.


Figure: KCL for the shown cut set yields $j_{1}(t)-j_{2}(t)+j_{3}(t)=0, \forall t$.

KCL equations

- originate from change conservation.
- are independent of the nature of the elements.
- are linear homogeneous equations with real coefficient $-1,0,1$.
- are dependent equations.


## KVL

## Loop

## Definition (Loop)

A subgraph of a graph is a loop if

- The subgraph is connected.
- Two branches of the subgraph are incident with each node of the subgraph.


Figure: Example of loop.

## Mesh and Super-mesh

## Statement (Mesh)

A mesh is a loop of a planar graph without any inner branch.

## Statement (Closed Chain)

A closed chain is a generalized loop of a planar graph that creates a closed path.


Figure: Examples of mesh and closed chain.

## KVL

## Definition (KVL)

For any lumped network and at any time, the algebraic sum of the aligned branch voltages around a loop (mesh, closed chain) is zero.


Figure: KVL for the shown loop yields $v_{4}(t)+v_{2}(t)-v_{5}(t)-v_{7}(t)+v_{8}(t)=0, \forall t$.

KVL equations

- originate from conservativity of electric field.
- are independent of the nature of the elements.
- are linear homogeneous equations with real constant coefficient $-1,0,1$.
- are dependent equations.


# Node-based Description 

## Number of Independent KCLs and Voltages

## Theorem (Number of Independent KCLs)

In a connected graph, the $n_{t}-1$ linear homogeneous algebraic equations obtained by applying KCL to each node except the reference node, constitute a set of linearly independent equations.

## Theorem (Number of Independent Voltages)

In a connected graph, the $n_{t}-1$ node voltages $\boldsymbol{e}$ measured with respect to the reference node constitute a set of linearly independent voltages.

## Node-to-branch Incidence Matrix

## Definition (Node-to-branch Incidence Matrix)

The node-to-branch incidence matrix $\boldsymbol{A}_{\boldsymbol{a}}$ is a rectangular matrix whose $(i, k)$ th element $a_{i k}$ is defined by

$$
a_{i k}= \begin{cases}1, & \text { if branch } k \text { leaves node } i \\ -1, & \text { if branch } k \text { enters node } i \\ 0, & \text { if branch } k \text { is not incident with node } i\end{cases}
$$

The matrix $\boldsymbol{A}_{\boldsymbol{a}}$ has dimension $n_{t} \times b$ and rank $n_{t}-1$, where $n_{t}$ and $b$ are the number of nodes and branches, respectively.

## Definition (Reduced Node-to-branch Incidence Matrix)

The reduced node-to-branch incidence matrix $\boldsymbol{A}$ is obtained from $\boldsymbol{A}_{\boldsymbol{a}}$ by eliminating the row corresponding to the reference node. The matrix $\boldsymbol{A}$ has dimension ( $n_{t}-$ 1) $\times b$ and is of full rank $n_{t}-1$.

## KCL Matrix Equation for Nodes

## Statement (KCL Matrix Equation for Nodes)

$\boldsymbol{A}_{a} \boldsymbol{j}=0$ describes $n_{t} K C L$ equations of the nodes, where $\boldsymbol{j}$ denotes branch currents vector.

## Statement (KCL Matrix Equation for Nodes)

$\boldsymbol{A} \boldsymbol{j}=0$ describes $n_{t}-1$ linearly independent KCL equations of the nodes.

## Statement (Branch Voltages)

The branch voltages $\boldsymbol{v}$ are obtained from the linearly-independent node voltages $\boldsymbol{e}$ by the equation $\boldsymbol{v}=\boldsymbol{A}^{T} \mathbf{e}$, where $\boldsymbol{A}^{T}$ is the transpose of $\boldsymbol{A}$.

## KCL Matrix Equation for Nodes

## Example (KCL Equation Matrix)

The circuit below has 3 independent KCL equations at its nodes.


$$
\boldsymbol{A}=\left[\begin{array}{ccccc}
1 & 1 & 0 & 0 & 0 \\
0 & -1 & 1 & 1 & 0 \\
0 & 0 & 0 & -1 & 1
\end{array}\right], \quad \boldsymbol{j}=\left[\begin{array}{l}
j_{1} \\
j_{2} \\
j_{3} \\
j_{4} \\
j_{5}
\end{array}\right], \quad \boldsymbol{A} \boldsymbol{j}=0, \quad\left\{\begin{array}{l}
j_{1}+j_{2}=0 \\
-j_{2}+j_{3}+j_{4}=0 \\
-j_{4}+j_{5}=0
\end{array}\right.
$$

## KCL Matrix Equation for Nodes

## Example (Branch Voltages)

The circuit below has 5 branch voltages.


$$
\boldsymbol{A}^{T}=\left[\begin{array}{ccc}
1 & 0 & 0 \\
1 & -1 & 0 \\
0 & 1 & 0 \\
0 & 1 & -1 \\
0 & 0 & 1
\end{array}\right], \quad \boldsymbol{e}=\left[\begin{array}{l}
e_{1} \\
e_{2} \\
e_{3}
\end{array}\right], \quad \boldsymbol{v}=\left[\begin{array}{l}
v_{1} \\
v_{2} \\
v_{3} \\
v_{4} \\
v_{5}
\end{array}\right], \quad \boldsymbol{v}=\boldsymbol{A}^{T} \boldsymbol{e}, \quad\left\{\begin{array}{l}
v_{1}=e_{1} \\
v_{2}=e_{1}-e_{2} \\
v_{3}=e_{2} \\
v_{4}=e_{2}-e_{3} \\
v_{5}=e_{3}
\end{array}\right.
$$

Mesh-based Description

## Topological Graphs

## Definition (Topological Graph)

Each different representation of a graph is called topological graph.


Figure: Three different topological graphs corresponding to a same graph. A loop remains unchanged for different topological graphs while a mesh may change.

## Planar Graphs

## Definition (Planar Graph)

A graph is planar if it can be drawn on the plane in such a way that no two branches intersect at a point which is not a node.

## Definition (Mesh and Outer-Mesh)

Any loop of a planar graph for which there is no branch in its interior is called a mesh. The loop of a planar graph for which there is no branch in its exterior is called the outer-mesh.


Figure: Examples of planar and non-planar graphs.

## Hinged Graphs

## Definition (Hinged Graph)

A graph is hinged if it can be partitioned into two non-isolated sub-graphs which are connected together by one node.


Figure: Examples of hinged and unhinged graphs. Circuit analysis of a hinged graph simplifies to separate analysis of its unhinged sub-graphs provided that there is no coupling between the unhinged sub-graphs.

## Number of Meshes

## Theorem (Number of Meshes)

For a connected unhinged planar graph, the number of meshes is equal $I=b-n_{t}+1$, where $b$ is the number of branches and $n_{t}$ is the number of nodes.


Figure: A planar unhinged graph with $n_{t}=9$ nodes, $b=14$ branches, and $I=14-9+1=6$ meshes.

## Number of Independent KVLs and Currents

## Theorem (Number of Independent KVLs)

In a connected planar unhinged graph, the $b-n_{t}+1$ linear homogeneous algebraic equations obtained by applying KVL to each mesh except the outer mesh constitute a set of linearly independent equations.

## Theorem (Number of Independent Currents)

In a connected planar unhinged graph, the $b-n_{t}+1$ mesh currents $\boldsymbol{i}$ constitute a set of linearly independent currents.

## Mesh-to-branch Incidence Matrix

## Definition (Mesh-to-branch Incidence Matrix)

The mesh-to-branch incidence matrix $\mathbf{M}_{\mathbf{a}}$ is a rectangular matrix whose $(i, k)$ th element $m_{i k}$ is defined by
$m_{i k}=$
1, if branch $k$ is in mesh or outer-mesh $i$ and their directions coincide
$\{-1$, if branch $k$ is in mesh or outer-mesh $i$ and their directions don't coincide 0 , if branch $k$ does not belong to mesh or outer-mesh $i$

The matrix $\boldsymbol{M}_{\mathbf{a}}$ has dimension $(I+1) \times b$ and rank $I$, where $I$ and $b$ are the number of meshes and branches, respectively.

## Definition (Reduced Mesh-to-branch Incidence Matrix)

The reduced mesh-to-branch incidence matrix $\boldsymbol{M}$ is obtained from $\boldsymbol{M}_{\mathbf{a}}$ by eliminating the row corresponding to the outer mesh. The matrix $\boldsymbol{M}$ has dimension $I \times b$ and is of full rank $I$.

## KVL Matrix Equation for Meshes

## Statement (KVL Matrix Equation for Meshes)

$\boldsymbol{M}_{\mathbf{a}} \boldsymbol{v}=0$ describes I +1 KVL equations of the meshes, where $\boldsymbol{v}$ denotes branch voltages vector.

## Statement (KVL Matrix Equation for Meshes)

$\mathbf{M v}=0$ describes I linearly independent KVL equations of the meshes.

## Statement (Branch Currents)

The branch currents $\boldsymbol{j}$ are obtained from the linearly-independent mesh currents $\boldsymbol{i}$ by the equation $\boldsymbol{j}=\boldsymbol{M}^{\top} \boldsymbol{i}$, where $\boldsymbol{M}^{\top}$ is the transpose of $\boldsymbol{M}$.

## KVL Matrix Equation for Meshes

## Example (KVL matrix equation)

The circuit below has 3 independent KVL equations at its meshes.


$$
\boldsymbol{M}=\left[\begin{array}{ccccc}
1 & 1 & 0 & 0 & 0 \\
0 & -1 & 1 & 1 & 0 \\
0 & 0 & 0 & -1 & 1
\end{array}\right], \quad \boldsymbol{v}=\left[\begin{array}{l}
v_{1} \\
v_{2} \\
v_{3} \\
v_{4} \\
v_{5}
\end{array}\right], \quad \boldsymbol{M} \boldsymbol{v}=0, \quad\left\{\begin{array}{l}
v_{1}+v_{2}=0 \\
-v_{2}+v_{3}+v_{4}=0 \\
-v_{4}+v_{5}=0
\end{array}\right.
$$

## KVL Matrix Equation for Meshes

## Example (Branch currents)

The circuit below has 5 branch currents.


$$
\boldsymbol{M}^{T}=\left[\begin{array}{ccc}
1 & 0 & 0 \\
1 & -1 & 0 \\
0 & 1 & 0 \\
0 & 1 & -1 \\
0 & 0 & 1
\end{array}\right], \quad \boldsymbol{i}=\left[\begin{array}{l}
i_{1} \\
i_{2} \\
i_{3}
\end{array}\right], \quad \boldsymbol{j}=\left[\begin{array}{l}
j_{1} \\
j_{2} \\
j_{3} \\
j_{4} \\
j_{5}
\end{array}\right], \quad \boldsymbol{j}=\boldsymbol{M}^{T} \boldsymbol{i}, \quad\left\{\begin{array}{l}
j_{1}=i_{1} \\
j_{2}=i_{1}-i_{2} \\
j_{3}=i_{2} \\
j_{4}=i_{2}-i_{3} \\
j_{5}=i_{3}
\end{array}\right.
$$

## Cut Set-based Description

## Trees

## Definition (Tree of a Connected Graph)

A graph is called the tree of a connected graph if
(1) It is a connected sub-graph.
(2) It contains all the nodes of the connected graph.

- It contains no loops.



## Trees

## Definition (Tree Branch)

The branches of a tree of a connected graph are called tree branch.

## Definition (link Branch)

The branches of a connected graph not in its associated tree are called link branch.


Graph $G$


Tree $T_{1}$


Tree $T_{2}$


Tree $T_{3}$


Tree $T_{4}$

Figure: Examples of trees of a graph.

## Trees

## Theorem (Fundamental Theory of Graphs)

Given a connected graph $G$ of $n_{t}$ nodes and $b$ branches, and a tree $T$ of $G$,

- There is a unique path along the tree between any pair of nodes.
- There are $n_{t}-1$ tree branches and $b-n_{t}+1$ links.
- Every link of $G$ and the unique tree path between its nodes constitute a unique loop (this is called the fundamental loop associated with the link).
- Every tree branch of $T$ together with some links defines a unique cut set. This cut set is called a fundamental cut set associated with the tree branch.


## Corollary (Fundamental Theory of Graphs)

Suppose that $G$ has $n_{t}$ nodes, $b$ branches, and s separate parts. Let $T_{l}, T_{2}, \cdots, T_{s}$ be trees of each separate part, respectively. The set $\left\{T_{l}, T_{2}, \cdots, T_{s}\right\}$ is called a forest of $G$. Then the forest has $n_{t}-s$ branches, $G$ has $b-n_{t}+s$ links, and the remaining statements of the fundamental theorem are true.

## Trees

## Example (Fundamental cut sets)

The circuit below has 4 fundamental cut sets. The direction of each cut set is inherited from the direction of its associated tree branch.


## Trees

## Example (Fundamental loops)

The circuit below has 4 fundamental loops. The direction of each loop is inherited from the direction of its associated link branch.


## Number of Independent KCLs and Voltages

## Theorem (Number of Independent KCLs)

In a connected graph, the $n_{t}-1$ linear homogeneous algebraic equations obtained by applying KCL to the fundamental cut sets of a tree of the graph, constitute a set of linearly independent equations.

## Theorem (Number of Independent Voltages)

In a connected graph, the $n_{t}-1$ tree branch voltages constitute a set of linearly independent voltages.

## Fundamental Cut Set Matrix

## Definition (Fundamental Cut Set Matrix)

The fundamental cut set matrix $\boldsymbol{Q}$ is a rectangular matrix whose $(i, k)$ th element $q_{i k}$ is defined by
$q_{i k}=$
(1, if branch $k$ belongs to cut set $i$ and has the same direction
$\{-1, \quad$ if branch $k$ belongs to cut set $i$ and has the opposite direction
0 , if branch $k$ does not belong to cut set $i$
The matrix $\boldsymbol{Q}$ has dimension $\left(n_{t}-1\right) \times b$ and is of full rank $n_{t}-1$, where $n_{t}-1$ and $b$ are the number of tree branches and branches, respectively.

## KCL Matrix Equation for Cut Sets

## Statement (KCL Matrix Equation for Cut Sets)

$\boldsymbol{Q} \boldsymbol{j}=0$ describes $n_{t}-1$ linearly independent KCL equations of the cut sets, where $j$ denotes branch currents vector.

## Statement (Branch Voltages)

The branch voltages $\boldsymbol{v}$ are obtained from the linearly-independent tree branch voltages $\boldsymbol{e}$ by the equation $\boldsymbol{v}=\boldsymbol{Q}^{T} \boldsymbol{e}$, where $\boldsymbol{Q}^{T}$ is the transpose of $\boldsymbol{Q}$.

## KCL Matrix Equation for Cut Sets

## Example (KCL matrix equation)

The circuit below has 4 independent KCL equations at its cut sets.

$\boldsymbol{Q}=\left[\begin{array}{cccccccc}1 & -1 & 0 & 0 & 1 & 0 & 0 & 0 \\ -1 & 1 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & -1 & -1 & -1 & 0 & 0 & 1 & 0 \\ 0 & -1 & -1 & 0 & 0 & 0 & 0 & 1\end{array}\right], \quad \boldsymbol{j}=\left[\begin{array}{l}j_{1} \\ j_{2} \\ j_{3} \\ j_{4} \\ j_{5} \\ j_{6} \\ j_{7} \\ j_{8}\end{array}\right], \quad \boldsymbol{Q} \boldsymbol{j}=0, \quad\left\{\begin{array}{l}j_{1}-j_{2}+j_{5}=0 \\ -j_{1}+j_{2}+j_{3}+j_{4}+j_{6}=0 \\ -j_{2}-j_{3}-j_{4}+j_{7}=0 \\ -j_{2}-j_{3}+j_{8}=0\end{array}\right.$

## KCL Matrix Equation for Cut Sets

## Example (Branch voltages)

The circuit below has 8 branch voltages.


$$
\boldsymbol{Q}^{T}=\left[\begin{array}{cccc}
1 & -1 & 0 & 0 \\
-1 & 1 & -1 & -1 \\
0 & 1 & -1 & -1 \\
0 & 1 & -1 & 0 \\
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1
\end{array}\right], \quad \boldsymbol{e}=\left[\begin{array}{l}
e_{1} \\
e_{2} \\
e_{3} \\
e_{4}
\end{array}\right], \quad \boldsymbol{v}=\left[\begin{array}{l}
v_{1} \\
v_{2} \\
v_{3} \\
v_{4} \\
v_{5} \\
v_{6} \\
v_{7} \\
v_{8}
\end{array}\right], \quad \boldsymbol{v}=\boldsymbol{Q}^{T} \boldsymbol{e}, \quad\left\{\begin{array}{l}
v_{1}=e_{1}-e_{2} \\
v_{2}=-e_{1}+e_{2}-e_{3}-e_{4} \\
v_{3}=e_{2}-e_{3}-e_{4} \\
v_{4}=e_{2}-e_{3} \\
v_{5}=e_{1} \\
v_{6}=e_{2} \\
v_{7}=e_{3} \\
v_{8}=e_{4}
\end{array}\right]
$$

## Loop-based Description

## Number of Independent KVLs and Currents

## Theorem (Number of Independent KVLs)

In a connected graph, the $I=b-n_{t}+1$ linear homogeneous algebraic equations obtained by applying KVL to the fundamental loops of a tree of the graph, constitute a set of linearly independent equations.

## Theorem (Number of Independent Currents)

In a connected graph, the $I=b-n_{t}+1$ link branch currents constitute a set of linearly independent voltages.

## Fundamental Loop Matrix

## Definition (Fundamental Loop Matrix)

The fundamental loop matrix $\boldsymbol{B}$ is a rectangular matrix whose $(i, k)$ th element $b_{i k}$ is defined by

$$
b_{i k}=
$$

1, if branch $k$ is in loop $i$ and their directions agree
$\{-1, \quad$ if branch $k$ is in loop $i$ and their directions don't agree $0, \quad$ if branch $k$ is not in loop $i$

The matrix $B$ has dimension $I \times b$ and is of full rank $I$, where $I$ and $b$ are the number of link branches and branches, respectively.

## KVL Matrix Equation for Loops

## Statement (KVL Matrix Equation for Loops)

$B \boldsymbol{v}=0$ describes $I=b-n_{t}+1$ linearly independent KVL equations of the loops, where $\boldsymbol{v}$ denotes branch voltages vector.

## Statement (Branch Current)

The branch currents $\boldsymbol{j}$ are obtained from the linearly-independent tree link currents $\boldsymbol{i}$ by the equation $\boldsymbol{j}=\boldsymbol{B}^{T} \boldsymbol{i}$, where $\boldsymbol{B}^{T}$ is the transpose of $\boldsymbol{B}$.

## KVL Matrix Equation for Loops

## Example (KVL matrix equation)

The circuit below has 4 independent KVL equations at its loops.


Links 1, 2, 3, 4
Tree branches 5, 6, 7, 8

$$
\boldsymbol{B}=\left[\begin{array}{cccccccc}
1 & 0 & 0 & 0 & -1 & 1 & 0 & 0 \\
0 & 1 & 0 & 0 & 1 & -1 & 1 & 1 \\
0 & 0 & 1 & 0 & 0 & -1 & 1 & 1 \\
0 & 0 & 0 & 1 & 0 & -1 & 1 & 0
\end{array}\right], \quad \boldsymbol{v}=\left[\begin{array}{l}
v_{2} \\
v_{3} \\
v_{4} \\
v_{5} \\
v_{6} \\
v_{7}
\end{array}\right], \quad \boldsymbol{B} \boldsymbol{v}=0, \quad\left\{\begin{array}{l}
v_{1}-v_{5}+v_{6}=0 \\
-v_{2}+v_{5}-v_{6}+v_{7}+v_{8}=0 \\
v_{3}-v_{6}+v_{7}+v_{8}=0 \\
v_{4}-v_{6}+v_{7}=0
\end{array}\right.
$$

## KVL Matrix Equation for Loops

## Example (Branch currents)

The circuit below has 8 branch currents.


$$
\boldsymbol{B}^{T}=\left[\begin{array}{cccc}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1 \\
-1 & 1 & 0 & 0 \\
1 & -1 & -1 & -1 \\
0 & 1 & 1 & 1 \\
0 & 1 & 1 & 0
\end{array}\right], \quad \boldsymbol{i}=\left[\begin{array}{l}
i_{1} \\
i_{2} \\
i_{3} \\
i_{4}
\end{array}\right], \quad \boldsymbol{j}=\left[\begin{array}{l}
j_{1} \\
j_{2} \\
j_{3} \\
j_{4} \\
j_{5} \\
j_{6} \\
j_{7} \\
j_{8}
\end{array}\right], \quad \boldsymbol{j}=\boldsymbol{B}^{T} \boldsymbol{i}, \quad\left\{\begin{array}{l}
j_{1}=i_{1} \\
j_{2}=i_{2} \\
j_{3}=i_{3} \\
j_{4}=i_{4} \\
j_{5}=-i_{1}+i_{2} \\
j_{6}=i_{1}-i_{2}-i_{3}-i_{4} \\
j_{7}=i_{2}+i_{3}+i_{4} \\
j_{8}=i_{2}+i_{3}
\end{array}\right.
$$

## Comparison of Different Descriptions

## Different Equations

- Node-based: $\begin{cases}\mathrm{KCL}: & \boldsymbol{A} \boldsymbol{j}=0 \\ \mathrm{KVL:} & \boldsymbol{v}=\boldsymbol{A}^{T} \boldsymbol{e}\end{cases}$
- Mesh-based: $\begin{cases}\mathrm{KVL}: & \boldsymbol{M v}=0 \\ \mathrm{KCL}: & \boldsymbol{j}=\boldsymbol{M}^{T} \boldsymbol{i}\end{cases}$
- Cut Set-based: $\begin{cases}\mathrm{KCL}: & \boldsymbol{Q} \boldsymbol{j}=0 \\ \mathrm{KVL}: & \boldsymbol{v}=\boldsymbol{Q}^{\boldsymbol{T}} \boldsymbol{e}\end{cases}$
- Loop-based: $\begin{cases}\mathrm{KVL}: & \boldsymbol{B} \boldsymbol{v}=0 \\ \mathrm{KCL}: & \boldsymbol{j}=\boldsymbol{B}^{\boldsymbol{T}} \boldsymbol{i}\end{cases}$


## Number of Items

- Number of linearly independent KCLs: $n=n_{t}-1$
- Number of linearly independent voltages: $n=n_{t}-1$
- Number of tree branches: $n=n_{t}-1$
- Number of linearly independent KVLs: $I=b-n_{t}+1$
- Number of linearly independent currents: $I=b-n_{t}+1$
- Number of link branches: $I=b-n_{t}+1$
- Number of trees: $\left|\boldsymbol{A} \boldsymbol{A}^{T}\right|=\left|\boldsymbol{M} \boldsymbol{M}^{T}\right|=\left|\boldsymbol{Q} \boldsymbol{Q}^{T}\right|=\left|\boldsymbol{B} \boldsymbol{B}^{T}\right|$
- Number of trees in complete graph: $n_{t}^{n_{t}-2}$


## Relationship of Matrices

## Statement (KVL Matrix Equation for Loops)

Call $\boldsymbol{B}$ the fundamental loop matrix and $\boldsymbol{Q}$ the fundamental cut-set matrix of the same directed graph G, and let both matrices pertain to the same tree $T$. Then, $\boldsymbol{B Q}^{T}=0$ and $\boldsymbol{Q B}^{T}=0$. Furthermore, if we number the links from 1 to $I$ and number the tree branches from $I+1$ to $b$, then $\boldsymbol{B}_{\mid \times b}=\left[\boldsymbol{I}_{1 \times 1} \mid \boldsymbol{F}\right]$ and $\boldsymbol{Q}_{\left(n_{t}-1\right) \times b}=\left[-\boldsymbol{F}^{\top} \mid \boldsymbol{I}_{\left(n_{t}-1\right) \times\left(n_{t}-1\right)}\right]$.

$$
\begin{gathered}
\boldsymbol{Q}=0 \Rightarrow \quad \boldsymbol{Q}\left(\boldsymbol{B}^{T} \boldsymbol{i}\right)=0 \Rightarrow\left(\boldsymbol{Q} \boldsymbol{B}^{\boldsymbol{T}}\right) \boldsymbol{i}=0 \Rightarrow \quad \boldsymbol{Q} \boldsymbol{B}^{T}=0 \Rightarrow \quad \boldsymbol{B} \boldsymbol{Q}^{T}=0 \\
\boldsymbol{B} \boldsymbol{Q}^{\boldsymbol{T}}=0 \Rightarrow\left[\begin{array}{lll}
\boldsymbol{I}_{\mid \times 1} & \mid & \boldsymbol{F}_{l \times\left(n_{t}-1\right)}
\end{array}\right]\left[\frac{\boldsymbol{E}_{1 \times\left(n_{t}-1\right)}^{T}}{\boldsymbol{I}_{\left(n_{t}-1\right) \times\left(n_{t}-1\right)}}\right]=\boldsymbol{E}_{l \times\left(n_{t}-1\right)}^{T}+\boldsymbol{F}_{l \times\left(n_{t}-1\right)}=0
\end{gathered}
$$

## Relationship of Matrices

## Example (Possible equality of $\boldsymbol{A}$ and $\boldsymbol{Q}$ )

There may be a special tree for which the node-to-branch incident matrix $\boldsymbol{A}$ and fundamental cut set matrix $\boldsymbol{Q}$ are the same.


$$
\boldsymbol{A}=\boldsymbol{Q}=\left[\begin{array}{cccccccc}
1 & -1 & 0 & 0 & 1 & 0 & 0 & 0 \\
-1 & 0 & 0 & 1 & 0 & 1 & 0 & 0 \\
0 & 0 & 1 & -1 & 0 & 0 & 1 & 0 \\
0 & 1 & -1 & 0 & 0 & 0 & 0 & 1
\end{array}\right]
$$

## Relationship of Matrices

## Example (Possible equality of $M$ and $B$ )

There may be a special tree for which the mesh-to-branch incident matrix $\boldsymbol{M}$ and fundamental loop matrix $\boldsymbol{B}$ are the same.


$$
\begin{aligned}
& \boldsymbol{A}=\boldsymbol{Q}=\left[\begin{array}{cccccccc}
1 & -1 & 0 & 0 & 1 & 0 & 0 & 0 \\
-1 & 0 & 0 & 1 & 0 & 1 & 0 & 0 \\
0 & 0 & 1 & -1 & 0 & 0 & 1 & 0 \\
0 & 1 & -1 & 0 & 0 & 0 & 0 & 1
\end{array}\right] \\
& \boldsymbol{M}=\boldsymbol{B}=\left[\begin{array}{cccccccc}
1 & 0 & 0 & 0 & -1 & 1 & 0 & 0 \\
0 & 1 & 0 & 0 & 1 & 0 & 0 & -1 \\
0 & 0 & 1 & 0 & 0 & 0 & -1 & 1 \\
0 & 0 & 0 & 1 & 0 & -1 & 1 & 0
\end{array}\right]
\end{aligned}
$$

## Relationship of Matrices

## Example (Desired set of independent voltages)

The shown tree corresponds to a set of independent voltages that includes $v_{2}$ and $v_{6}$ and does not include $v_{1}, v_{3}$, and $v_{7}$.


## Relationship of Matrices

## Example (Desired set of independent currents)

The shown tree corresponds to a set of independent currents that includes $j_{4}$ and $j_{6}$ and does not include $j_{2}$ and $j_{7}$.


## Duality

## Dual Graphs

## Statement (Dual Graphs)

Two connected, unhinged, and planar topological graphs $G$ and $\hat{G}$ are dual if,

- There is a one-to-one correspondence between the meshes of $G$ (including the outer mesh) and the nodes of $\hat{G}$.
- There is a one-to-one correspondence between the meshes of $\hat{G}$ (including the outer mesh) and the nodes of $G$.
- There is a one-to-one correspondence between the branches of each graph in such a way that whenever two meshes of one graph have the corresponding branch in common, the corresponding nodes of the other graph have the corresponding branch connecting these nodes.


## Dual Graphs

## Example (Dual Graphs)

The two graphs below are dual of each other.

(3)



## Dual Graphs

## Example (Dual Directed Graphs)

The two directed graphs below are dual of each other.
(4)



## Dual Circuits

## Statement (Dual Circuits)

Two circuits are dual if,

- Their associative graphs, $G$ and $\hat{G}$, are dual.
- The governing circuit equations of $\hat{G}$ are obtained by the following replacement from governing circuit equations of $G$.

$$
\begin{aligned}
j & \rightarrow \hat{v} \\
v & \rightarrow \hat{j} \\
q & \rightarrow \hat{\phi} \\
\phi & \rightarrow \hat{q}
\end{aligned}
$$

## Dual Circuits

| $G$ | $\hat{G}$ |
| :---: | :---: |
| KVL | KCL |
| KCL | KVL |
| Node | Mesh |
| Mesh | Node |
| Refrence Node | Outer Mesh |
| Outer Mesh | Reference Node |
| Parallel Connection | Series Connection |
| Series Connection | Parallel Connection |
| Link Branch | Tree Branch |
| Tree Branch | Link Branch |
| Open Circuit | Short Circuit |
| Short Circuit | Open Circuit |



Table: Dual items.


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## Dual Circuits

## Example (Dual Graphs)

Two circuits below are dual.


$$
\frac{-1}{j \omega L} E_{1}+\left(j \omega C_{2}+G+\frac{1}{j \omega L}\right) E_{2}=0, \quad \frac{-1}{j \omega \hat{C}} \hat{l}_{1}+\left(j \omega \hat{L}_{2}+\hat{R}+\frac{1}{j \omega \hat{C}}\right) \hat{l}_{2}=0
$$

## Tellegen's Theorem

## Tellegen's Theorem

## Theorem (Tellegen's Theorem)

Consider an arbitrary lumped network whose graph $G$ has $b$ branches and $n_{t}$ nodes. Suppose that to each branch of the graph we assign arbitrarily a branch voltage $v_{k}$ and a branch current $j_{k}$ for $k=1,2 \cdots, b$, and suppose that they are measured with respect to arbitrarily picked associated reference directions. If the branch voltages $v_{1}, v_{2}, \cdots, v_{b}$ satisfy all the constraints imposed by KVL and if the branch currents $j_{1}, j_{2}, \cdots, j_{b}$, satisfy all the constraints imposed by $K C L$, then

$$
\sum_{k=1}^{b} v_{k} j_{k}=0
$$

## Tellegen's Theorem

- If the voltage sets $\left\{v_{k} \mid k=1, \cdots, b\right\}$ and $\left\{\hat{v}_{k} \mid k=1, \cdots, b\right\}$ and the current sets $\left\{j_{k} \mid k=1, \cdots, b\right\}$ and $\left\{\hat{j}_{k} \mid k=1, \cdots, b\right\}$ satisfy KVL and KCL requirements, then

$$
\sum_{k=1}^{b} v_{k} j_{k}=0, \quad \sum_{k=1}^{b} \hat{v}_{k} j_{k}=0, \quad \sum_{k=1}^{b} v_{k} \hat{j}_{k}=0, \quad \sum_{k=1}^{b} \hat{v}_{k} \hat{j}_{k}=0
$$

- Tellegen's theorem is independent of the nature of elements.
- Instantaneous and apparent power conservation are special cases of Tellegen's theorem.


## Tellegen's Theorem

## Example (Two-measurement experiment)

For the LTI RLC network below, Tellegen's theorem forces $\hat{J}_{1}=J_{2}$ in the two illustrated measurement scenarios.


$$
V_{1} \hat{\jmath}_{1}+V_{2} \hat{\jmath}_{2}+\sum_{k=3}^{b} V_{k} \hat{\jmath}_{k}=V_{1} \hat{\jmath}_{1}+V_{2} \hat{\jmath}_{2}+\sum_{k=3}^{b} z_{k} J_{k} \hat{\jmath}_{k}=0
$$

$$
\hat{V}_{1} J_{1}+\hat{V}_{2} J_{2}+\sum_{k=3}^{b} \hat{V}_{k} J_{k}=\hat{V}_{1} J_{1}+\hat{V}_{2} J_{2}++\sum_{k=3}^{b} z_{K} \hat{J}_{k} J_{k}=0
$$

$$
V_{1} \hat{J}_{1}+V_{2} \hat{J}_{2}=\hat{V}_{1} J_{1}+\hat{V}_{2} J_{2}
$$

$$
V_{s} \hat{\jmath}_{1}=V_{s} J_{2}
$$

$$
\hat{J}_{1}=J_{2}
$$

## Driving-point Impedance



Figure: Driving-point impedance of a passive RLCMT network. Coupled inductors are replaced with their passive equivalent circuits. Tranformers do not consume power.

- Complex power conservation: $-0.5 V_{1} J_{1}^{*}+0.5 \sum_{k=2}^{b} V_{k} J_{k}^{*}=0$
- Complex power conservation:

$$
0.5 Z_{i n}(j \omega)\left|J_{1}\right|^{2}=0.5 \sum_{R} R_{k}\left|J_{k}\right|^{2}+0.5 j \omega \sum_{L} L_{k}\left|J_{k}\right|^{2}-0.5 j \omega^{-1} \sum_{C} C_{k}^{-1}\left|J_{k}\right|^{2}
$$

- Driving-point impedance:

$$
Z_{i n}(j \omega)=\frac{\sum_{R} R_{k}\left|J_{k}\right|^{2}}{\left|J_{1}\right|^{2}}+j \frac{\sum_{L} \omega L_{k}\left|J_{k}\right|^{2}-\sum_{C} \omega^{-1} c_{k}^{-1}\left|J_{k}\right|^{2}}{\left|J_{1}\right|^{2}}=\Re\left\{Z_{i n}(j \omega)\right\}+j \Im\left\{Z_{i n}(j \omega)\right\}
$$

- Passivity condition: $\Re\left\{Z_{i n}(j \omega)\right\} \geq 0, \quad \Im\left\{Z_{i n}(j \omega)\right\} \in \mathbb{R}, \quad\left|\angle Z_{i n}(j \omega)\right| \leq \pi / 2$


## Driving-point Impedance



Figure: Driving-point impedance of a passive RLCMT network. Coupled inductors are replaced with their passive equivalent circuits. Tranformers do not consume power.

- Average dissipated power: $P_{a v_{k}}=0.5 R_{k}\left|J_{k}\right|^{2}$
- Average stored magnetic energy: $\overline{\mathcal{L}}_{L_{k}}=0.25 L_{k}\left|J_{k}\right|^{2}$
- Average stored electrical energy: $\overline{\mathcal{E}}_{c_{k}}=0.25 C_{k}\left|V_{k}\right|^{2}=0.25 C_{k}^{-1} \omega^{-2}\left|J_{k}\right|^{2}$
- Complex power conservation:
$0.5 Z_{i n}(j \omega)\left|J_{1}\right|^{2}=0.5 \sum_{R} R_{k}\left|J_{k}\right|^{2}+0.5 j \omega \sum_{L} L_{k}\left|J_{k}\right|^{2}-0.5 j \omega^{-1} \sum_{C} C_{k}^{-1}\left|J_{k}\right|^{2}$
- Complex power conservation:
$S=\sum_{R} P_{\mathrm{av}_{k}}+2 j \omega\left(\sum_{L} \overline{\mathcal{E}}_{L_{k}}-\sum_{C} \overline{\mathcal{E}}_{C_{k}}\right)=P_{\mathrm{av}}+2 j \omega\left(\overline{\mathcal{E}}_{L}-\overline{\mathcal{E}}_{C}\right)$
- Driving-point impedance: $z_{i n}(j \omega)=\frac{2 P_{a v}+4 j \omega\left(\bar{\varepsilon}_{L}-\bar{\varepsilon}_{c}\right)}{\left|N_{1}\right|^{2}}$


## Driving-point Impedance

| Circuit Type | $\Re\left\{Z_{\text {in }}(j \omega)\right\}$ | $\Im\left\{Z_{\text {in }}(j \omega)\right\}$ | $Z_{\text {in }}(j \omega)$ |
| :---: | :---: | :---: | :---: |
| RT | $\Re\left\{Z_{\text {in }}(j \omega)\right\} \geq 0$ | $\Im\left\{Z_{\text {in }}(j \omega)\right\}=0$ | $\angle Z_{\text {in }}(j \omega)=0$ |
| RLMT | $\Re\left\{Z_{\text {in }}(j \omega)\right\} \geq 0$ | $\Im\left\{Z_{\text {in }}(j \omega)\right\} \geq 0$ | $0 \leq \angle Z_{\text {in }}(j \omega) \leq \pi / 2$ |
| RCT | $\Re\left\{Z_{\text {in }}(j \omega)\right\} \geq 0$ | $\Im\left\{Z_{\text {in }}(j \omega)\right\} \leq 0$ | $-\pi / 2 \leq \not Z_{\text {in }}(j \omega) \leq 0$ |
| RLCMT | $\Re\left\{Z_{\text {in }}(j \omega)\right\} \geq 0$ | $\Im\left\{Z_{\text {in }}(j \omega)\right\} \in \mathbb{R}$ | $-\pi / 2 \leq / Z_{\text {in }}(j \omega) \leq \pi / 2$ |
| LCMT | $\Re\left\{Z_{\text {in }}(j \omega)\right\}=0$ | $\Im\left\{Z_{\text {in }}(j \omega)\right\} \in \mathbb{R}$ | $\angle Z_{\text {in }}(j \omega)= \pm \pi / 2$ |
| LMT | $\Re\left\{Z_{\text {in }}(j \omega)\right\}=0$ | $\Im\left\{Z_{\text {in }}(j \omega)\right\} \geq 0$ | $\angle Z_{\text {in }}(j \omega)=\pi / 2$ |
| CT | $\Re\left\{Z_{\text {in }}(j \omega)\right\}=0$ | $\Im\left\{Z_{\text {in }}(j \omega)\right\} \leq 0$ | $\angle Z_{\text {in }}(j \omega)=-\pi / 2$ |

Table: Driving-point impedance of passive networks.

## The End

